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# Space saving calculation of symbolic resultants

Jane M Pearson and Noel G Lloyd

**Abstract.** We describe an approach to the computation of symbolic resultants in which factors are removed during the course of the calculation, so reducing the stack size required for intermediate expressions and the storage space needed. We apply the technique to three well-established methods for calculating resultants. We demonstrate the advantages of our approach when the resultants are large and show that some otherwise intractable problems can be resolved. In certain cases a significant reduction in the cpu time required to calculate the resultant is also evident.

**Keywords.** symbolic computation, multivariate resultants.

## 1. Introduction

Our interest in the calculation of symbolic resultants arose from our research into some of the properties of systems of nonlinear differential equations (see for example [9],[10],[11],[12]). In order to provide a context for our discussion we first give a brief description of the mathematical background. We consider differential systems of the form

$$\dot{x} = \lambda x + y + p(x, y), \quad \dot{y} = -x + \lambda y + q(x, y), \quad (1)$$

where  $p$  and  $q$  are polynomials without linear or constant terms. The corresponding complex form of system (1) is

$$i\dot{z} = (1 + i\lambda)z + \sum_{i+j=2} A_{ij}z^i\bar{z}^j, \quad (2)$$

where  $z = x + iy$ ,  $A_{ij} \in \mathbb{C}$ . When  $\lambda = 0$  the origin is said to be a *fine focus*. Our objective is to establish the conditions under which the origin is a centre and to determine the maximum number of limit cycles that can be bifurcated from the origin for systems of the form (1) or (2) under perturbation of the coefficients. **All** orbits in the neighbourhood of the origin when it is a *centre* are closed; in contrast a *limit cycle* is an isolated closed orbit.

We separate the calculation of the conditions for the origin to be a centre into two parts: necessity and sufficiency. It is in the calculation of the necessary conditions that most of the large resultant calculations arise. These conditions are obtained by calculating the *focal values*, which are polynomials in the coefficients in  $p$  and  $q$ , (or in the  $A_{ij}$ ) and are defined as follows. There is a function  $V$ , analytic in a neighbourhood of the origin, such that the rate of change along orbits,  $\dot{V} = \eta_2 r^2 + \eta_4 r^4 + \dots$ , where  $r^2 = x^2 + y^2$  and  $\eta_2 = \lambda$ . The  $\eta_{2k}$  are the focal values and the number of terms in each  $\eta_{2k}$  grows rapidly as  $k$  increases. In examples where  $\eta_4$  has only four terms it is not unusual for  $\eta_{12}$  to have over a thousand terms, and  $\eta_{14}$  over three thousand.

The origin is a centre if, and only if, all the focal values are zero. By the Hilbert basis theorem, the ideal they generate has a finite basis so there is  $M$  such that if  $\eta_{2j} = 0$ , for  $j \leq M$ , then  $\eta_{2j} = 0$  for all  $j$ . The value of  $M$  is not known *a priori*. The origin is a fine focus of order  $k$  if  $\eta_{2m} = 0$ , for  $m \leq k$  but  $\eta_{2k+2} \neq 0$ . At most  $k$  small amplitude limit cycles can bifurcate from a fine focus of order  $k$ .

Our approach is to calculate the first few focal values for a given system and to make substitutions from each one into the other focal values. We have  $\eta_2 = \lambda = 0$ , necessarily. We make a substitution from  $\eta_4 = 0$ , for one of the variables, into subsequent calculated focal values. Then we make substitutions from  $\eta_6 = 0$ ,  $\eta_8 = 0$  and so on. We thus obtain expressions for each of the eliminated variables in terms of the remaining variables - information that is required for the bifurcation of the limit cycles. After each substitution we remove common factors from the remaining calculated focal values, these being candidates for the conditions under which the origin is a centre. We continue until we can show that if the remaining factor of focal value  $\eta_{2k}$  is zero then focal value  $\eta_{2k+2}$  is necessarily non-zero. Then the maximum order of the origin as a fine focus is  $k$ . The sufficiency of the candidate centre conditions is proved independently using a range of different techniques.

We have considered looking for a Gröbner basis for the set of focal values; there are three main drawbacks to this approach. First, we do not know *a priori* the value of  $M$  for a given differential system. Secondly, the Gröbner basis does not readily give us the information we require in order to bifurcate the limit cycles. Finally, obtaining the Gröbner basis is non-trivial for many systems.

In the systems of interest to us the focal values usually involve at least seven variables. As each variable is eliminated the remaining focal values grow; they contain more terms, the variables occur to higher degrees and the integer coefficients become larger. At each stage of the elimination process an attempt is made to simplify the focal values by factorising them, each such factor is then considered individually. However we inevitably reach a point where the variable we wish to eliminate does not occur linearly in any of the focal values (or factors of the focal values) and we must employ polynomial remainder sequences, as in [10],[12] or use resultant calculations, see for example [12], to eliminate that variable. Often successive resultant calculations are required; the performance of such calculations is sensitive to the order in which the variables are eliminated.

Typically we find the variable being eliminated occurs to degree greater than 10, with the other variables occurring to higher degrees. The integer coefficients often have as many digits as the total degree of the polynomial. Procedures for calculating resultants fail for some examples simply because the expressions involved become so massive. Techniques to make the computations more feasible are essential. In this paper we describe how we have developed software to calculate the resultant of multivariate polynomials which exploits the particular features we have observed in previous examples. Many of the resultants we wish to calculate have several simple factors occurring to high multiplicity; we aim to remove as many of these factors as we can during the course of the calculation, so reducing the size of intermediate expressions. Some of these factors, such as common factors of leading coefficients of the variable being eliminated, can be predicted. Others are determined by factorisations or greatest common divisor (gcd) calculations.

The resultant can be calculated in various ways; we consider three approaches, each of which requires the evaluation of the determinant of a certain matrix. The Sylvester, Bézout and Companion matrices which we use are described in section 4. The most time and space consuming element of any of these approaches is the calculation of the determinant, the setting up of the matrix always being a minor consideration. In section 5 we present a method for calculating the determinant in which intermediate expression swell is reduced by removing factors of the resultant as they arise. This idea can be applied to any method that involves the calculation of a determinant.

There are techniques to compute the resultant of three polynomials with respect to two variables directly, using for example the Macaulay determinant [14],[15] or Dixon resultants [8]. The Dixon resultant is a generalisation of the Bézout-Cayley method which is described below. A major limitation of the Dixon based resultants is that often an extraneous factor is generated, a problem addressed in [2]. The size of the matrix is dependent on the variable ordering and it is essential to minimize this to enable the effective calculation of the determinant. The size of the Macaulay matrix is given by  $\binom{d-2}{2}$ , where  $d$  is the sum of the total degree of the two variables being eliminated in each of the three polynomials. In the examples we encounter  $d$  is likely to be at least 50. Although we do not consider this approach here it would be possible to apply the idea of the early removal of factors to these resultant calculations.

We demonstrate our approach by reference to two examples which come from the investigation of two systems which we describe in section 3. Section 2 contains a general discussion of resultant calculations. In section 6 we present two further examples which highlight the advantages of the Bézout or Companion matrix approaches. Our concluding remarks are in section 7.

## 2. Resultants

Typically we have two multivariate polynomials, say  $f, g$ , each with several hundred, and often several thousand, terms and we wish to establish under what circumstances  $f = g = 0$ . We denote the resultant of  $f$  and  $g$  with respect to the variable  $x$  by  $\text{res}(f, g, x)$ . We have  $f = g = 0$  only if  $\text{res}(f, g, x) = 0$ . To simplify further calculations we require all the irreducible factors of  $\text{res}(f, g, x)$ .

The computation of symbolic resultants of large multivariate polynomials (for example polynomials of degree greater than ten in the variable being eliminated with coefficients that are polynomials in one, or more, variables occurring to degrees greater than ten) is very demanding of both computer space and time. It is often impossible to obtain the resultants we want using the currently available software. Of the techniques available most involve the calculation of the determinant of a matrix, the exact form of the matrix reflecting the different methods. It is also possible to calculate resultants using interpolation techniques [13]. First the degree of the required resultant is determined, then resultant calculations are performed for specific values of one of the variables. Finally a polynomial is interpolated from this data set. For this to be effective one requires an efficient means of calculating the resultants and an efficient method for interpolating polynomials in several variables.

We note that in many of our examples the resultant contains many simple, repeated factors and we exploit this by removing any such factors as they arise during the calculation of the determinant. In this paper we concentrate on the calculation of the resultant with respect to one variable but our approach could easily be applied to any method which requires a determinant to be computed. Our computations were performed on a Compaq Alpha XP1000 workstation, with single 667MHz Alpha EV5 processor and 1 Gb of memory, using the Computer Algebra systems REDUCE and Maple.

The resultant of two given polynomials with respect to a given variable can be thought of as the elimination of the given variable from the two polynomials. Where the polynomials have a non-trivial gcd their resultant vanishes. Writing the irreducible, multivariate polynomials  $f$  and  $g$  as polynomials in the single variable  $x$ , with polynomial coefficients in the remaining variables, we have

$$f = \sum_{i=0}^n a_i x^i \text{ and } g = \sum_{i=0}^m b_i x^i, \quad (3)$$

where  $m \geq n$ . Our requirement is to find all factors of  $\text{res}(f, g, x)$ . We present a technique for identifying some of these factors during the course of the calculation of the resultant. Removing such factors during the computation reduces the size of the intermediate expressions calculated and consequently makes the overall calculation that much more feasible.

### 3. Examples

We compare three different methods, based on the Sylvester, Bézout and Companion matrices respectively, by reference to examples. The first two examples are straightforward, we include them merely to demonstrate the differences in the three methods that we consider. In particular we are interested in when factors of the resultant first occur in the calculation. Two further examples are presented in section 6, these illustrate the effectiveness of our technique.

**Example 1** The first example arose in the investigation of cubic differential systems of the form

$$i\dot{z} = z + A_{11}z\bar{z} + A_{30}z^3 + A_{12}z\bar{z}^2 + A_{03}\bar{z}^3, \quad (4)$$

where  $A_{ij} \in \mathbb{C}$ . At a certain point in the elimination of variables from the focal values for this system we require  $R_0 + aR_1 = 0$ , where  $R_0(b, c, d), R_1(b, c, d)$  are real non-homogeneous polynomials of degrees 5, 4 respectively in  $d$ , and  $a, b, c, d$  are functions of the real and imaginary parts of the  $A_{ij}$ . We need to consider the two possibilities  $R_1 \neq 0, a = -\frac{R_0}{R_1}$  and  $R_0 = R_1 = 0$ . Here we are concerned with the second case and we calculate the resultant of  $R_0$  and  $R_1$  with respect to  $d$ . We refer to this as Example 1.

**Example 2** The second example arose in the investigation of differential systems of the form

$$\begin{aligned} \dot{x} &= y(1 + kx), \\ \dot{y} &= -x + c_1x^2 + c_2xy + c_3y^2 + c_4x^3 + c_5x^2y + c_6xy^2 + c_7y^3, \end{aligned} \quad (5)$$

when (5) has two coexisting fine foci. The origin is a fine focus and we can scale the system so that there is a second fine focus at  $(1, 0)$ ; hence  $c_4 = 1 - c_1, c_5 = -c_2$  and  $(k+1)(c_1-2) > 0$ . The variables  $k, c_6, c_7$  are eliminated from the focal values leaving polynomials in the remaining variables  $c_1, c_2, c_3$ . Further details can be found in [7]. In particular we have two polynomials  $S_0, S_1$  of degrees 6, 4 in  $c_2^2$  with coefficients that are polynomials in  $c_1$  and  $c_3$ . We refer to the resultant of  $S_0$  and  $S_1$  with respect to  $c_2^2$  as Example 2.

### 4. The matrices

Considering the polynomials in (3) we first define the form of the matrices used in the three approaches to the calculation of resultants and note some of their properties.

#### 4.1. Sylvester matrix

The Sylvester matrix,  $S$ , is defined as

$$S = \begin{pmatrix} b_m & b_{m-1} & \cdot & \cdot & \cdot & b_0 & 0 & \cdot & \cdot & 0 \\ 0 & b_m & b_{m-1} & \cdot & \cdot & b_1 & b_0 & 0 & \cdot & 0 \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & \cdot & \cdot & 0 & b_m & b_{m-1} & \cdot & \cdot & \cdot & b_0 \\ a_n & a_{n-1} & \cdot & \cdot & a_0 & 0 & \cdot & \cdot & \cdot & 0 \\ 0 & a_n & a_{n-1} & \cdot & \cdot & a_0 & 0 & \cdot & \cdot & 0 \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & \cdot & \cdot & \cdot & 0 & a_n & \cdot & \cdot & \cdot & a_0 \end{pmatrix} \begin{matrix} \text{row 1} \\ \cdot \\ \cdot \\ \cdot \\ n \\ n+1 \\ \cdot \\ \cdot \\ \cdot \\ n+m \end{matrix}$$

where the first  $n$  rows of the matrix contain only the coefficients of the polynomial  $g$ , followed by  $m$  rows with only the coefficients of  $f$ . The Sylvester matrix is generally large, its order being the sum of the degrees of the polynomials  $f$  and  $g$ , in  $x$ , but it is sparse in that at most  $2mn + m + n$  of its elements are non-zero. The Sylvester matrix need not be stored explicitly for our procedure to determine the resultant, since we require only the individual coefficients  $a_i$  and  $b_i$ . The resultant of  $f$  and  $g$  is given by

$$\text{res}(f, g, x) = (-1)^{mn} \det(S);$$

as is explained in [16] and discussed further in [3]. The sign of the resultant is not important for our purposes as we are only interested in the conditions under which the resultant is zero. Clearly common factors of the coefficients occurring in any column of the Sylvester matrix are factors of the resultant. In practice we find that, in our problems, only columns 1 and  $m + n$  are likely to have such factors.

#### 4.2. Bézout matrix

The approach to finding a resultant using the Bézout matrix is often known as Cayley's method [1],[6]. Let

$$\beta(x, y) = \frac{f(x)g(y) - f(y)g(x)}{(x - y)}.$$

Clearly  $(x - y)$  is a factor of the numerator, so  $\beta$  is a polynomial of degree  $m - 1$  in  $x$  and  $y$ . The elements of the Bézout matrix,  $B$ , are given by

$$B(i, j) = \text{coefficient of } x^{i-1}y^{j-1} \text{ in } \beta$$

for  $i, j = 1, 2, \dots, m$ . The resultant in this case is given by

$$\text{res}(f, g, x) = \pm b_m^{n-m} \det(B),$$

where  $b_m \neq 0$ . The  $m \times m$  Bézout matrix is symmetric, since  $\beta(x, y) = \beta(y, x)$ , so it is only necessary to store the upper (or lower) triangle of  $B$ . We consider those

elements in the upper triangle, which has the following form:

$$\begin{pmatrix} B_{1,1} & B_{1,2} & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & B_{1,m} \\ & B_{2,2} & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & B_{2,m} \\ & & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ & & & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ & & & & B_{n+1,n+1} & \cdot & \cdot & \cdot & \cdot & B_{n+1,m} \\ & & & & & B_{n+2,n+2} & \cdot & \cdot & B_{n+2,m-1} & 0 \\ & & & & & & \cdot & \cdot & 0 & 0 \\ & & & & & & & \cdot & \cdot & \cdot \\ & & & & & & & & \cdot & 0 \end{pmatrix}$$

The entries can be determined according to a procedure described in [5] suitably modified for  $m \neq n$ . Let  $ab_{i,j} = a_j b_{i-1} - a_{i-1} b_j$ , with  $a_p = 0$  if  $p > n$ . Then  $B(1,j) = ab_{1,j}$  for  $j = 1, \dots, m$ ,  $B(i,m) = ab_{i,m}$  for  $i = 2, \dots, \min(m, n+1)$ , and  $B(i,j) = ab_{i,j} + B(i-1, j+1)$  for  $i = 2, \dots, m$ ,  $j = i, \dots, m-1$ . The number of elements to be stored is  $\tau = \frac{1}{2}m(m+1)$  when  $m = n$ ,  $\tau - \frac{1}{4}(m-n-1)(m-n+1)$  when  $m-n > 0$  is odd and  $\tau - \frac{1}{4}(m-n-2)(m-n+2) - 1$  when  $m-n \geq 2$  is even.

The expression for  $\det(B)$  differs from the resultant being sought by a factor of  $b_m^{m-n}$ , if  $m \neq n$ . We must remove at least the  $m-n$  factors  $b_m$  during the course of the resultant calculation to avoid creating an expression which is even larger than the required resultant.

**Lemma 1.** *If  $m-n=1$  then  $b_m$  is a factor of row, or column,  $m$ . If  $m-n > 1$  then  $b_m$  is a factor of row and column  $m$ .*

*Proof.* We have  $B(i,m) = B(m,i) = -a_{i-1}b_m$ ,  $i = 1, \dots, n+1$  and if  $m-n > 1$  then  $B(i,m) = B(m,i) = 0$ ,  $i = n+2, \dots, m$ . The proof follows.  $\square$

We note that as we are working with only the upper triangle of  $B$  the factor  $b_m$  can be removed from row/column  $m$  only if  $m-n > 1$ .

### 4.3. Companion matrix

The Companion matrix for the polynomial  $f$  is defined as

$$C_f = \begin{pmatrix} 0 & \cdot & \cdot & \cdot & -a_0/a_n \\ 1 & 0 & \cdot & \cdot & -a_1/a_n \\ 0 & 1 & 0 & \cdot & -a_2/a_n \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & \cdot & \cdot & 1 & -a_{n-1}/a_n \end{pmatrix}$$

where the leading coefficient of  $f$  is strictly non-zero. We note that  $C_f$  is an  $n \times n$  matrix, the lowest ordered matrix of the three being considered, and we have, as described in [4], that

$$\text{res}(f, g, x) = \pm a_n^m \det(g(C_f)).$$

The form of the matrix  $C_f$  is such that  $C_f^2$  is the matrix  $C_f$  with column one removed and column  $n$  with entries  $A_{1,1}/a_n^2$ ,  $A_{1,2}/a_n^2$ ,  $\dots$ ,  $A_{1,n}/a_n^2$ , where  $A_{1,1} =$



$a_0 a_{n-1}$  and  $A_{1,i} = a_{i-1} a_{n-1} - a_{i-2} a_n$ , for  $i = 2, \dots, n$ . In general  $C_f^p$  is  $C_f^{p-1}$  with column one removed and a new column  $n$  with entries  $A_{p-1,i}/a_n^p$ , where  $A_{0,i} = -a_{i-1}$ ,  $A_{k,1} = -a_0 A_{k-1,n}$ ,  $A_{k,i} = a_n A_{k-1,i-1} - a_{i-1} A_{k-1,n}$  for  $k = 2, \dots, m-1$  and  $i = 1, \dots, n$ . Hence the elements of the matrix  $G = g(C_f)$  are given by

$$G(i, j) = \left( \sum_{t=n-j+1}^m A_{t-n+j-1,i} b_t a_n^{m-t} + \Omega \right) / a_n^{m-n+j}, \quad (6)$$

where  $\Omega = 0$  if  $i < j$  and  $\Omega = b_{i-j} a_n^{m-n+j}$  if  $i \geq j$ , for  $i, j = 1, \dots, n$ . The matrix  $G$  is dense. We note that the elements of column  $j$  of the matrix  $G$  have a common factor of  $a_n^{-(m-n+j)}$ . Writing  $C$  as the matrix  $G$  with these common factors removed we have

$$\text{res}(f, g, x) = \pm \prod_{i=m-n+1}^{m-1} a_n^{-i} \det(C) = \pm a_n^{(n-1)(\frac{1}{2}n-m)} \det(C).$$

We refer to  $C$  as the reduced Companion matrix for  $f$  and  $g$ .

**Lemma 2.** *The elements of  $C$  have a common factor that is the greatest common divisor of  $a_n$  and  $b_m$ .*

*Proof.* From (6) every element of  $C$  has  $a_n$  as a factor unless  $t = m$  in which case  $b_m$  is a factor.  $\square$

As  $m \geq n$  the determinant we need to calculate is a larger expression than the required resultant, therefore it is essential to develop a technique in which the factors  $a_n$  are removed as they arise.

## 5. The determinant

The methods described above require the evaluation of the determinant of a symbolic matrix to obtain the resultant. There are efficient methods for calculating the determinant of a numerical matrix and most can be applied to matrices with symbolic entries. However the time taken when the matrix contains symbolic elements can be significantly longer. This is partly because mathematical operations on symbolic expressions are performed by software rather than hardware.

Suppose that  $M = (M_{ij})$  is an  $s \times s$  matrix. The determinant of  $M$  is usually defined recursively, as the sum of certain matrix elements times the determinants of matrices that are the cofactors of those elements. Although conceptually simple, computationally this is not a good definition to use as the time taken to evaluate the determinant of a matrix  $M$  grows exponentially as  $s$  increases. We adopt a direct approach in which we start by evaluating the determinants of  $2 \times 2$  matrices and progress through the determinants of  $3 \times 3$  matrices until we eventually calculate  $\det(M)$ . This allows us to remove common factors of the sub-determinants at each stage of the process.

Let  $H_k(i_1, i_2, \dots, i_k)$  be the  $k \times k$  sub-determinant formed from rows  $i_1, \dots, i_k$  of the first  $k$  columns of  $M$ , with  $H_1(i) = -M_{i,1}$ . We define iteratively

$$H_k(i_1, i_2, \dots, i_k) = \sum_{r=1}^k (-1)^{r-1} M_{i_r, k} H_{k-1}(i_1, \dots, i_{r-1}, i_{r+1}, \dots, i_k) \quad (7)$$

for  $k = 2, \dots, s$ . Then  $\det(M) = H_s(1, \dots, s)$ . For each  $k$  there are potentially  $\binom{s}{k}$  sub-determinants to be calculated.

**Lemma 3.** *Any common factor of the  $k \times k$  sub-determinants is also a factor of  $\det(M)$ .*

*Proof.* If  $\alpha$  is a common factor of the  $H_k$  for  $k = \kappa$  then by (7) it is also a factor of the  $H_{\kappa+1}$ . It follows that  $\alpha$  is a factor of  $\det(M) = H_s$ .  $\square$

The approach we adopt is to remove common factors from the  $H_k$ , before calculating the  $H_{k+1}$ . In general the number and size of the expressions for the  $H_k$  is such that calculating greatest common divisors (or factorizing each one) at each level is non-trivial. So we remove any common factors that we know will be present, for certain  $H_k$ , before attempting a gcd calculation. For example, in the Bézout and reduced Companion matrices we know that the leading coefficient  $b_m$  and multiplicities of  $a_n$  respectively must be present as factors in certain of the  $H_k$ . In the Sylvester matrix we know that for  $k < n$  only common factors of  $a_n$  and  $b_m$  can arise in the  $H_k$ . Factors of the leading coefficients occur repeatedly in the resultant, so we check for them being common factors in the  $H_k$  even when we do not explicitly predict their presence. Integer factors of the  $H_k$  are easily lifted and a greatest common divisor for them can be readily computed. Removing this common factor in itself stems the growth of the integer coefficients. Other common factors that may occur depend very much on the problem and our objective is to develop a strategy of looking for such factors at the earliest opportunity.

### 5.1. Determinant of the Sylvester matrix

The calculation of the determinant of the Sylvester matrix using our method gives rise to some unique problems due to the sparsity of the matrix. It is desirable not to waste time calculating any  $H_k$  that is known *a priori* to be zero and, more importantly, we do not wish to calculate non-zero  $H_k$  that are not required in the determination of  $H_{k+1}$ . We have  $H_1(1) = -b_m$ ,  $H_1(n+1) = -a_n$  and all other  $H_1$  are zero. There are five non-zero  $H_2$ , namely

$$\begin{aligned} H_2(1, 2) &= b_m^2, & H_2(1, n+1) &= b_m a_{n-1} - b_{m-1} a_n, \\ H_2(1, n+2) &= b_m a_n, & H_2(2, n+1) &= -H_2(1, n+2), & H_2(n+1, n+2) &= a_n^2. \end{aligned}$$

When  $k \leq n$ , each  $k \times k$  sub-matrix can notionally be reduced to an upper triangular matrix, the determinant of which is zero if there is a zero entry on the diagonal. Referring to (7) we say that  $i_1, i_2, \dots, i_k$  are the parameters of  $H_k$ . Let  $h_{k,i} = \{i_1, i_2, \dots, i_k\}$  for each  $H_k$ , where  $i = 1, \dots, \text{number of } H_k$ . The determinants we must calculate for  $k = \kappa \leq n$  are as follows:

- the  $H_\kappa$  with parameters  $h_{\kappa-1,i}$  plus either  $\kappa$  or  $n + \kappa$ ,
- the  $H_\kappa$  for matrices formed from  $\kappa$  of the rows occupied in column  $\kappa - 1$ . We define a row to be of type  $r$  if the first entry in the row is in column  $r$ . The rows chosen must be of types 1 to  $\kappa$ , two selected rows of type  $\rho$  generate a row of type  $\rho + 1$ . The selection of  $\kappa$  rows from the possible  $2(\kappa - 1)$  rows made up of two rows of each type  $r$ , for  $r = 1, \dots, \kappa - 1$ , will include at least two rows of the same type. We refer to the selection of two rows of the same type, that is rows  $i$  and  $n + i$ , as an ‘and’, as opposed to an ‘or’ when row  $i$  or row  $n + i$  is chosen. We consider all possible combinations of ‘and’s and ‘or’s - if  $p$  ‘and’s are included then  $\kappa - 2p$  ‘or’s are required, for  $p = 1, \dots, \lfloor \frac{\kappa}{2} \rfloor$ . We ensure that our selection results in one row of each of the required types, otherwise the sub-determinant is zero. For example, row 1 or row  $n + 1$ , or both must be chosen and no  $h_{\kappa,i}$  can involve rows  $\kappa - 1$ ,  $n + \kappa - 1$ ,  $\kappa - 2$ , and  $n + \kappa - 2$  simultaneously.

For  $k = n + 1, \dots, m - 1$  we could modify the technique used for  $k \leq n$ ; for  $k = \kappa$  add  $n + \kappa$  to each  $h_{\kappa,i}$  and consider combinations of rows occupied in column  $\kappa - 1$ . However this leads to a small number of sub-determinants being calculated that are not required. We adopt an alternative approach of recursively deleting from the integers  $Z = \{1, 2, \dots, m + n\}$  the occupied rows of columns  $m + n, m + n - 1, \dots, \kappa + 1$  to determine the  $h_{\kappa,i}$ . Starting with  $Z$  we first delete either  $n$  or  $m + n$ , then from  $Z - \{n\}$  we delete  $n - 1$  or  $m + n - 1$  or  $m + n$  and from  $Z - \{m + n\}$  we delete  $n - 1$  or  $m + n - 1$ , and so on. A few of the sub-determinants calculated will be zero but this is considered preferable to calculating non-zero quantities that are not needed.

The form of the Sylvester matrix means that we can determine which  $H_k$  are required for  $k = m, \dots, m + n - 1$  from a knowledge of which  $H_j$ , for  $j = 1, \dots, n$  are to be calculated. For each of the parameters in  $h_{j,i} = \{i_1, i_2, \dots, i_j\}$ , if  $i_r \leq n$  then  $j_r = m + n - i_r + 1$  else  $j_r = 2n - i_r + 1$  is deleted from the integers  $1, 2, \dots, m + n$  to give  $h_{m+n-j,i}$ . For example, assuming that  $n \geq 2$ , the five non-zero  $H_{m+n-2}$  are

$$\begin{aligned} &H_{m+n-2}(1, \dots, m + n - 2), H_{m+n-2}(1, \dots, n - 2, n + 1, \dots, m + n), \\ &H_{m+n-2}(1, \dots, n - 2, n, \dots, m + n - 1), \\ &H_{m+n-2}(1, \dots, n - 1, n + 1, \dots, m + n - 1), \\ &H_{m+n-2}(1, \dots, n - 1, n + 1, \dots, m + n - 2, m + n). \end{aligned}$$

The number of sub-determinants which must be calculated rises steeply as  $k$  increases, reaching a maximum when  $k = \lfloor \frac{m+n}{2} \rfloor$ . We have already seen that five sub-determinants are required when  $k = 2$ . The following table shows the number required for values of  $k \leq 10 \leq n$ . The increase in number as  $k$  goes from  $n + 1$  to  $\lfloor \frac{m+n}{2} \rfloor$  is not as rapid and for every value of  $k$  significantly fewer sub-determinants are required than would be the case for a dense matrix.

$k$	2	3	4	5	6	7	8	9	10
no. of $H_k$	5	14	42	132	429	1430	4862	16796	58786

Having decided which of the sub-determinants are to be calculated we turn to the question of common factors and their occurrence.

**Lemma 4.** *Common factors of the sub-determinants  $H_k$ , for  $k \leq n$ , for the Sylvester matrix are common factors of the leading coefficients of the polynomials  $f, g$ . Common factors of the  $H_k$ , for  $n+1 < k \leq m$  are factors of the leading coefficient of  $f$ .*

*Proof.* We have that when  $k \leq n$ ,  $H_k(1, 2, \dots, k) = b_m^k$  and when  $k \leq m$ ,  $H_k(n+1, n+2, \dots, n+k) = a_n^k$ . The proof follows.  $\square$

We cannot expect to find any other common factors of the  $H_k$  for  $k < m+1$ , so we defer the calculation of greatest common divisors until we reach values of  $k \geq m+1$ . At this point the number of sub-determinants for each  $k$  is decreasing, although the expressions generated for each one are large.

### Sylvester algorithm

*Begin*

*Input  $f, g$*

$H_1(1) = b_m / \gcd(a_n, b_m)$

$H_1(n+1) = a_n / \gcd(a_n, b_m)$

$H_2(1, 2) = H_1(1)b_m$

$H_2(1, n+1) = H_1(1)a_{n-1} - b_{m-1}H_1(n+1)$

$H_2(1, n+2) = H_1(n+1)b_m$

$H_2(n+1, n+2) = H_1(n+1)a_n$

*Remove factors of the  $\gcd(a_n, b_m)$  that are common to the  $H_2$*

$H_2(2, n+1) = -H_2(1, n+2)$

$k = 3$  *step 1 until  $n$*

*Determine the  $h_{k,i}$  for each  $H_k$  to be calculated*

*Add  $k$  or  $n+k$  to  $h_{k-1,i}$*

*Select  $k$  occupied rows of column  $k-1$ , of different row types*

*Calculate the  $H_k$*

*Remove common numerical factor from the  $H_k$*

*Remove factors of  $a_n$  that are common to the  $H_k$*

$Z = \{1, 2, \dots, m+n\}$

$k = n+1$  *step 1 until  $m-1$*

*Determine the  $h_{k,i}$  for each  $H_k$  to be calculated*

*Delete from  $Z$  the occupied rows of columns  $m+n, m+n-1, \dots, k+1$*

*Calculate the  $H_k$*

*Remove common numerical factor from the  $H_k$*

*Remove factors of  $a_n$  that are common to the  $H_k$*

$k = m$  *step 1 until  $m+n-1$*

*Determine the  $h_{k,i}$  for each  $H_k$  to be calculated*

$r = 1$  *step 1 until  $m+n-k$*

$i_r = r^{\text{th}}$  *element of  $h_{m+n-k,i}$*

*If  $i_r \leq n$  then delete  $m+n-i_r+1$  else  $2n-i_r+1$  from  $Z$*

*Calculate the  $H_k$*   
*Remove common numerical factor from the  $H_k$*   
*Remove factors of  $a_n$  that are common to the  $H_k$*   
*If  $k > m$*   
*Find and remove common factors of the  $H_k$*   
*Calculate  $H_{m+n}$*   
*Return  $H_{m+n}$ , list of factors removed together with their multiplicities*  
*End of algorithm*

We illustrate this approach using the two examples described in section 3.

**Example 1a** We use our technique based on sub-determinants of the Sylvester matrix to calculate the resultant of the polynomials of Example 1, namely  $R_0$  and  $R_1$ , with respect to  $d$ . The degrees of  $d$  in  $R_0, R_1$  are 5, 4 respectively, so the resultant is given by  $H_9$ . The leading coefficients of  $d$  in  $R_0, R_1$  have the factor  $75(b+3)(3b-1)$  in common and  $(162b+5)\varphi_4(b)$ , where  $\varphi_4$  is a polynomial of degree 4 in  $b$ , is a factor of the leading coefficient in  $R_1$ . We find

$$\text{res}(R_0, R_1, d) = 16A(162b+5)^2\gamma^3\varphi_4^2(b)\varphi_3^2(b)\Psi_1(b, c)\Psi_2(b, c), \quad (8)$$

with  $\gamma = (3b^2c+12b^2+8bc+32b-3c-27)$ ,  $A = 1080000(3b-1)^3(b+3)(162b+5)\gamma$ ,  $\psi_3$  is a polynomial of degree 3 in  $b$ ,  $\Psi_1$  is a polynomial of degrees 8, 2 in  $b, c$  respectively and  $\Psi_2$  is degree 40 in  $b$ , 8 in  $c$ . Here  $A$  is removed from the resultant in the form of common factors of the  $H_k$  during the course of the procedure and the other factors are determined by factorisation of  $H_9$ . We remove  $75(b+3)(3b-1)$  from the  $H_1$ , then  $15(3b-1)$  is common to the  $H_2$  and the  $H_3$ . Subsequently only common numerical factors are encountered until the  $H_8$  are calculated, when  $(162b+5)\gamma$  is found.

**Example 2a** In Example 2 the polynomials  $S_0, S_1$  are of degrees 6, 4 in  $c_2^2$  with coefficients that are polynomials in  $c_1$  and  $c_3$ . We calculate the resultant

$$\text{res}(S_0, S_1, c_2^2) = 11520A(c_1 + c_3 - 1)^3(3c_1 + 3c_3 - 2)^2\Omega(c_1, c_3), \quad (9)$$

where  $\Omega(c_1, c_3)$  is of degree 44 in  $c_1, c_3$  and  $A = 4976640(c_1 + c_3 - 1)^5\delta$ , with  $\delta = (15c_1^2 + 30c_1c_3 - 20c_1 + 15c_3^2 - 20c_3 + 8)$ , is removed from the resultant in the form of common factors of the  $H_k$  during the course of the procedure. The other factors are obtained by factorising  $H_{10}$ . Here the common factor of the leading coefficients  $10(c_1 + c_3 - 1)\delta$  is removed from the  $H_1$ , then  $2(c_1 + c_3 - 1)$  is common to the  $H_2$ . Again only numerical common factors are found until the  $H_8$ , which have  $72(c_1 + c_3 - 1)$  in common. The remaining factor  $288(c_1 + c_3 - 1)^2$  of  $A$  is found in the  $H_9$ .

## 5.2. Determinant of the Bézout matrix

When applying our method of sub-determinant calculations to the Bézout matrix it is advantageous to re-order the columns so that any zero elements occur in the first few columns. This ensures that the  $H_k$  for  $k \leq m - n - 1$  are as simple as possible. Now the  $H_2(i, j)$  are non-zero for  $i = 1, 2, \dots, n+1 \leq m-1$ ,  $j = 2, \dots, n+2 \leq m$  and in general we must evaluate  $\binom{\min(m, n+k)}{k}$  sub-determinants for each  $k = 2, \dots, m$ . Symmetry allows us to store only the upper triangle of the Bézout matrix, but we must adjust our calculation of the  $H_k$  accordingly. By Lemma 1 if  $m - n > 1$  a factor  $b_m^2$  can be removed from the matrix, the  $m - n - 2$  factors  $b_m$  that remain to be removed to give the final expression for the determinant occur in the  $H_k$  for  $k = 2, \dots, m - n - 1$ . When  $m - n = 1$  then  $b_m$  is a factor of the  $H_2$ . Although the number of sub-determinant calculations required is considerably fewer than in the calculation of the determinant of the Sylvester matrix each individual sub-determinant is more complicated.

**Bézout algorithm***Begin**Input  $f, g$* *Store factors of  $b_m$  and their multiplicities**Set up upper triangular  $m \times m$  Bézout matrix* *$i = 1$  step 1 until  $\min(m, n + 1)$*  *$j = i$  step 1 until  $n$* 

$$B(i, j) = b_{i-1}a_j - b_ja_{i-1}$$

*if  $m > n$  then* *$j = n + 1$  step 1 until  $m$* 

$$B(i, j) = -a_{i-1}b_j$$

 *$i = 2$  step 1 until  $\min(m - 1, n + 1)$*  *$j = i$  step 1 until  $m - 1$* 

$$B(i, j) = B(i, j) + B(i - 1, j + 1)$$

 *$i = n + 2$  step 1 until  $(m + n + 1)/2$*  *$j = i$  step 1 until  $m + n - i + 1$* 

$$B(i, j) = B(i - 1, j + 1)$$

*If  $m - n > 1$  remove common factor  $b_m$  from column  $m$* *Calculate the  $\binom{\min(m, n+2)}{2}$  sub-determinants  $H_2$*  *$i = 1$  step 1 until  $m - 2$*  *$j = i + 1$  step 1 until  $m - 1$* 

$$H_2(i, j) = B(i, m - 1)B(j, m) - B(i, m)B(j, m - 1)$$

 *$i = 1$  step 1 until  $m - 1$* 

$$H_2(i, m) = B(i, m - 1)B(m, m) - B(i, m)B(m - 1, m)$$

*Find and remove common numerical factor of the  $H_2$* *Remove common non-numerical factor of  $b_m$  from the  $H_2$*  *$k = 3$  step 1 until  $m - 1$* *Combine column  $m - k + 1$  and the  $H_{k-1}$  to give the  $\binom{n}{k}H_k$* *Remove common numerical factor from the  $H_k$* *If  $k < m - n - 1$* *Remove non-numerical factors of  $b_m$  that are common to the  $H_k$* *Else**Find and remove common factors of the  $H_k$* *Calculate  $H_m$* *Return  $H_m$ , list of factors removed together with their multiplicities**End of algorithm*

**Example 1b** We calculate the resultant for the polynomials given in Example 1 using our technique applied to the Bézout matrix. Here  $m = 5, n = 4$ , so we cannot remove  $b_5$  from the stored matrix elements. Common factors of the  $H_k$  can occur for any value of  $k$ . Again we obtain the resultant as given by (8) with  $A$  removed during the course of the procedure. The common factor of the  $H_2$  is  $150b_5$ , where  $b_5 = 225(3b - 1)(b + 3)\varphi_5(b)$  and  $\varphi_5(b)$  is a polynomial of degree 5 in  $b$ . We have

$60(3b-1)^2(b+3)$  common to the  $H_3$  and  $120(3b-1)(162b+5)\gamma$  common to the  $H_4$ . The remainder of the determinant is given by the factorisation of  $H_5$ .

**Example 2b** Similarly we calculate the resultant for Example 2 and obtain the result given by (9). In this case  $b_6^2$  is removed from the matrix before the determinant is calculated, hence the resultant is equal to  $H_6$ . We find  $10(c_1+c_3-1)\delta$  is common to the  $H_2$ ,  $24(c_1+c_3-1)$  to the  $H_3$ ,  $72(c_1+c_3-1)$  to the  $H_4$  and finally  $288(c_1+c_3-1)^2$  to the  $H_5$ .

### 5.3. Determinant of the Companion matrix

We apply our method of sub-determinant calculations to the approach based on the Companion matrix. The number of sub-determinants  $H_k$  that must be calculated is  $\binom{n}{k}$  for  $k = 2, \dots, n-1$ . We note that although the matrix is dense fewer sub-determinants are required than for the Bézout or Sylvester matrices. The entries in the reduced Companion matrix are even more complex than those of the Bézout matrix.

**Lemma 5.** *The sub-determinants  $H_k$ , for fixed  $k > 2$ , for the reduced Companion matrix for  $f$  and  $g$  have a common factor  $a_n^{m-n+k-1}$  after a factor  $a_n^{m-n+k-2}$  has been removed from  $H_{k-1}$ .*

*Proof.* Consider first the  $H_2$ . We have for  $p = 1, 2, \dots, n-1$  and  $s = p+1, \dots, n$

$$\begin{aligned} H_2(p, s) &= C(p, 1)C(s, 2) - C(s, 1)C(p, 2) \\ &= \left( \sum_{t=n}^m A_{t-n,p} b_t a_n^{m-t} + b_{p-1} a_n^{m-n+1} \right) \left( \sum_{t=n-1}^m A_{t-n+1,s} b_t a_n^{m-t} + \Omega_s \right) \\ &\quad - \left( \sum_{t=n}^m A_{t-n,s} b_t a_n^{m-t} + b_{s-1} a_n^{m-n+1} \right) \left( \sum_{t=n-1}^m A_{t-n+1,p} b_t a_n^{m-t} + \Omega_p \right) \end{aligned}$$

where  $\Omega_v = 0$ , if  $v < 2$  and  $\Omega_v = b_{v-2} a_n^{m-n+2}$  otherwise. The coefficient of the term in  $b_m^2$  is

$$\begin{aligned} &A_{m-n,p} A_{m-n+1,s} - A_{m-n,s} A_{m-n+1,p} \\ &= (a_n A_{m-n-1,p-1} - a_{p-1} A_{m-n-1,n}) (a_n A_{m-n,s-1} - a_{s-1} A_{m-n,n}) \\ &\quad - (a_n A_{m-n-1,s-1} - a_{s-1} A_{m-n-1,n}) (a_n A_{m-n,p-1} - a_{p-1} A_{m-n,n}). \end{aligned}$$

Clearly  $a_n$  is a factor; repeated application of the recurrence relation for the  $A$ s gives  $a_n^{m-n+1}$  as a factor of the coefficient of  $b_m^2$ . Similarly for all the coefficients of  $b_i b_j$ . We conclude that  $a_n^{m-n+1}$  is a factor of the  $H_2$ , which is removed.

In general the  $H_k$  involve the entries of column  $k$  of the matrix  $C$ , for which the summations are over  $t = n-k+1, \dots, m$ . Terms in  $H_k$  independent of  $a_n$  cancel and repeated application of the recurrence relation for the  $A$ s gives  $a_n^{m-n+k-1}$  as a factor.  $\square$

According to Lemma 2 the gcd of  $a_n, b_m$  is a factor of each element of the matrix  $C$ , we remove this factor before calculating the sub-determinants. We also



remove any numerical factors that are common to the rows or columns of  $C$ . The common factors of the  $H_k$ , that we know will arise according to Lemma 5, are removed for each  $k$ . The result given in Lemma 5 means that  $H_n$  is divisible by  $a_n^{m-1}$ , and unless common factors of the  $H_k$ , for  $k < n$ , ultimately reduce  $H_n$  considerably the expression generated could be larger than that for the resultant.

### Companion algorithm

*Begin*

*Input  $f, g$*

*Calculate  $\gcd(a_n, b_m)$*

*Set up  $n \times m$  matrix  $A$*

*$j=1$  step 1 until  $n$*

$$A(0, j) = -a_{j-1}$$

$$A(1, 1) = A(0, 1)A(0, n)$$

*$i=2$  step 1 until  $m-1$*

$$A(i, 1) = A(0, 1)A(i-1, n)$$

*$j=2$  step 1 until  $n$*

$$A(i, j) = a_n A(i-1, j-1) + A(0, j)A(i-1, n)$$

*Set up  $n \times n$  reduced Companion matrix  $C$*

*$i=1$  step 1 until  $n$*

*$j=1$  step 1 until  $n$*

$$s=0$$

*$k=1$  step 1 until  $m$*

$$\text{If } k-n+j \geq 1 \text{ then } s = s + A(k-n+j-1, i)b_k a_n^{m-k}$$

$$\text{If } i \geq j \text{ then } s = s + a_n^{j+m-n} b_{i-j}$$

$$C(i, j) = s$$

*Remove  $\gcd(a_n, b_m)$  from each element of  $C$*

*Remove common numerical factor from each row of  $C$*

*Remove common numerical factor from each column of  $C$*

*Calculate the  $\binom{n}{2}$  sub-determinants  $H_2$*

*$i=1$  step 1 until  $n-1$*

*$j=i+1$  step 1 until  $n-1$*

$$H_2(i, j) = C(i, 1)C(j, 2) - C(i, 2)C(j, 1)$$

*Remove common numerical factor of the  $H_2$*

*Remove common factor  $a_n^{m-n+1} / \gcd(a_n, b_m)$  from the  $H_2$*

*Remove  $\gcd(a_n, b_m)^{\min(\max(0, m-2n+1), m-n+1)}$  from the  $H_2$*

*Find and remove any other common factors of the  $H_2$*

*$k=3$  step 1 until  $n$*

*Combine column  $k$  and the  $H_{k-1}$  to give the  $\binom{n}{k}$  sub-determinants  $H_k$*

*Remove common numerical factor of the  $H_k$*

*Remove common factor  $a_n^{m-n+k-1} / \gcd(a_n, b_m)$  from the  $H_k$*

*Remove  $\gcd(a_n, b_m)^{\min(\max(0, \frac{1}{2}(k-1)(2m+k)-kn), m-n+k-1)}$  from the  $H_k$*

*Find and remove any other common factors of the  $H_k$*

*End of algorithm*

**Example 1c** Again using the polynomials of Example 1 we calculate their resultant using the technique of sub-determinants for the Companion matrix. Common factor  $(3b - 1)(b + 3)$  of the leading coefficients is removed from the matrix  $C$ ; this is equivalent to extracting  $(3b - 1)^4(b + 3)^4$  from the determinant of  $C$ . Common numerical factors of  $75^9$ , the constant multiple in  $a_4^9$  (which the determinant must be divided by to give the resultant) and 720, which is a factor of the resultant, are also removed from  $C$ . The  $H_2$  with  $a_4^2$  removed have a common factor  $2(b + 3)(3b - 1)^3$ , and  $20(b + 3)(3b - 1)(162b + 5)\gamma$  is common to the  $H_3$  with  $a_4^3$  removed. Then  $H_4$ , before  $a_4^4$  is removed, is a polynomial of degree 96 in  $b$  whereas the resultant is of degree 77 in  $b$ . Clearly our objective of reducing the size of intermediate expressions has not been achieved for this example.

We find a similar drawback when evaluating the resultant of the polynomials for Example 2. In general the approach based on the Companion matrix leads to expressions that are larger than the required resultant. However, we note that this approach may be advantageous when the degree of  $f$  is small compared with that of  $g$ , as is demonstrated by Example 3 in the next section.

## 6. Further examples

We have shown how our approach can be used in conjunction with any of the three methods mentioned for the calculation of resultants. All three methods detect the same factors of the determinant during the course of its calculation, albeit at different stages in the procedures. We use whichever method is found to be the most effective for the case at hand, remembering that the motivation is to minimise the space required rather than the time taken. Although, in general, the Companion matrix approach requires more intermediate storage we find it is the most appropriate method for the first example in this section. Finally we present an example for which our method of sub-determinant calculations based on the Bézout matrix outperforms the resultant procedures within REDUCE and Maple, either in terms of space required or the time taken to perform the calculation.

**Example 3** This example serves to illustrate the advantage of the Companion matrix approach when one of the polynomials is of much lower degree than the other and, incidentally, highlights the size of the integer coefficients that occur. It is again based on polynomials which arose in the investigation of the differential system (4) but at the stage when only two variables  $a, b$  remain. Note that the first polynomial is homogeneous apart from the second term. If both polynomials were homogeneous and irreducible, then they would be simultaneously zero if and only if  $a = b = 0$ . In the course of the calculation we need to calculate the resultant of the polynomials  $\Upsilon, \Gamma$ , where

$$\begin{aligned} \Upsilon = & 464557320a^8 - 1633110051a^7 - 2037974188a^6b^2 - 21579017205a^5b^3 + \\ & 57307963738a^4b^4 + 189853264347a^3b^5 - 18545338008a^2b^6 + \\ & 30841802109ab^7 + 281875666338b^8, \end{aligned}$$

$$\begin{aligned}
\Gamma = & 8259534896175565488964a^{22} - 284138308132047340861208a^{21}b + \\
& 3708406249776621132785379a^{20}b^2 - 17017146098630322328970258a^{19}b^3 - \\
& 91500335842788491862840869a^{18}b^4 + 1532634666360967054531403566a^{17}b^5 - \\
& 6750066429361290171643370686a^{16}b^6 - 1231720038428668010567454736a^{15}b^7 \\
& + 117539646928826438620238727628a^{14}b^8 \\
& - 335865049700120874407070128512a^{13}b^9 \\
& - 476104610799723924571817647894a^{12}b^{10} \\
& + 4139408084167000893605158117052a^{11}b^{11} \\
& - 4108395371133928750343042289054a^{10}b^{12} \\
& - 18344976447322762629935442611460a^9b^{13} \\
& + 43386332084253057824995735934408a^8b^{14} \\
& + 21061401447247049003613260482560a^7b^{15} \\
& - 137067367296123029900260781375400a^6b^{16} \\
& + 35157616804635363171364106398536a^5b^{17} \\
& + 194641929399070622857185124352379a^4b^{18} \\
& - 63612019502265141475620194539674a^3b^{19} \\
& - 161119480899094799287341029858469a^2b^{20} \\
& - 245553624919122584912713549530ab^{21} \\
& + 97196908650313686371963778379950b^{22}.
\end{aligned}$$

We compared the calculation of  $\text{res}(\Upsilon, \Gamma, a)$  using various methods. First using the two procedures provided within REDUCE 3.7; the one which is based on the Bézout technique failed with “stack overflow” and the other, which uses sub-resultant calculations, obtained the resultant in 51s cpu. Our approach applied to the Bézout matrix fails as all space is exhausted and, because of the limitation on the dimension of arrays, our implementation of the Sylvester matrix approach cannot be used. Of the three routines described in this paper we can only obtain the resultant using the Companion matrix approach. This method is suited to this example because the degree of the variable being eliminated is much smaller in  $\Upsilon$  than it is in  $\Gamma$ . The resultant is obtained in 31s cpu, which compares favourably with the procedure provided as part of the REDUCE package, particularly as our code was designed to minimise the amount of space required not the time taken.

**Example 4** Consider the differential system

$$\begin{aligned}
\dot{x} &= x(x - fy + f - 1)(bx + y + c + d - b), \\
\dot{y} &= y(ax - y - a + 1)(dx + y + c - e).
\end{aligned} \tag{10}$$

This is an example of a Kolmogorov system; such systems are often used by ecologists to model the interaction of two species[12]. There are four points that can

be fine foci for system (10) but at most three of these can coexist. Here (10) has been scaled so that one of the fine foci is at  $(1, 1)$ , two other fine foci are  $S = (1 - m/(a+b), 1 - am/(a+b))$  and  $T = (1 + f(e-m)/(df+1), 1 + (e-m)/(df+1))$ , where  $m = c+d+1$ . We can transform (10) so that it is of the form (1) with origin at each of these points in turn. For each transformed system we must have  $\lambda = 0$ ; for the system with origin at  $(1, 1)$  this means  $e = 0$  and for the other two systems this condition gives expressions for the variables  $a, b$ . When considering the focal value  $\eta_4$  for the system with origin at  $S$  we need to investigate the possibility of  $P_1 = P_2 = 0$ . Here  $P_1$ , a polynomial with 389 terms, is degree 11 in each of the three remaining variables  $d, f, m$  and  $P_2$ , with 263 terms, is degree 9 in  $d, m$ , and degree 10 in  $f$ . Very few of the integer coefficients are greater than 1000. See the Appendix for details of  $P_1, P_2$ . Attempting to calculate  $\text{res}(P_1, P_2, m)$  using the REDUCE procedure based on the Bézout technique, or using the Maple resultant procedure, leads to all available space being exhausted. The resultant can be computed using the REDUCE procedure which uses sub-resultant calculations giving  $\text{res}(P_1, P_2, m) = 512d^{23}f^{48}\Phi$ , where  $\Phi$  is a polynomial of degrees 56 in  $d$ , 138 in  $f$ . This calculation took over 33 hours cpu time, of which more than 67% was taken by garbage collection. Using our procedure based on the Bézout matrix we find, during the calculation of the determinant, that  $256d^{14}f^{37}(f-1)^{58}(f+1)^2(df+1)^{32}$  is a factor of the resultant with cofactor  $2d^9f^{11}\Theta$ , where  $\Theta$  is a polynomial of degrees 24 in  $d$ , 46 in  $f$ . This calculation required approximately 13.5 hours cpu time with less than 3% of the time being spent in garbage collection. Also the factorisation of  $\Theta$ , to obtain the resultant in its simplest form, is trivial compared with that of  $\Phi$ . There is a clear advantage in using our technique when the resultant has many simple factors to high multiplicity.

We have

$$\begin{aligned} \text{res}(P_1, P_2, m) = & 512d^{23}(d-1)f^{48}(f-1)^{76}(f+1)^4(d+f)(df+1)^{49} \times \\ & (d-f^2+f+1)^2(d^2f^2-2df^2+3df+f^2-2f+1) \times \\ & (df^2+2df+d-2f^2+4f+2). \end{aligned}$$

This result leads to the conclusion that the three coexisting fine foci can each be of order at most one.

## 7. Conclusion

We have presented three different approaches to the calculation of the resultant of two multivariate polynomials and have compared their performance with reference to examples. The technique we use differs from other methods in that some factors of the resultant are removed during the course of its calculation. This reduces the amount of storage space required and the stack size needed for individual expressions. Obtaining all the factors of the resultant in their simplest terms is expedited as the factorisation of one very large expression is replaced by several factorisations of simpler functions.

The approach based on the Sylvester matrix seems to offer the least potential. It requires the calculation of by far the most sub-determinants, although each sub-determinant in the early stages contains fewer terms than those for the other methods. It also has the added complication of the need to determine which sub-determinants to calculate. For efficiency we store the parameters  $h_{k,i}$  in an array rather than a list in our implementation using REDUCE. There is a maximum array size set within the REDUCE system we use and we soon reach this limit. Currently we cannot consider polynomials of degree greater than nine in the variable being eliminated using this software. In the Sylvester matrix approach no common factors other than factors of the leading coefficients can arise in the  $H_k$ , for  $k \leq m$ , so if  $m - n$  is small there is little opportunity to detect more interesting factors of the resultant at an early stage in the calculation of the  $H_k$ . Conversely, when using the Bézout matrix technique with  $m - n \leq 2$  it is possible that such factors will arise from  $H_3$  onwards.

In certain examples expressions generated during the Companion matrix approach are even larger than the resultant being sought, but we have given one instance, in Example 3, where this method outperforms the other two and those procedures which are provided as part of the REDUCE 3.7 package. When the degree of the variable being eliminated is much lower in one of the polynomials than the other the Companion matrix approach can be advantageous. In general the Bézout matrix algorithm is the most useful. Example 4, where the resultant has many simple factors of high multiplicity, clearly demonstrates the advantage of removing factors of the determinant during its calculation.

The extent of any advantage given by the removal of factors during the course of the calculation is very much problem dependent. We have found that our method, based on the calculation of sub-determinants and the removal of factors of the resultant as they arise, offers advantages over other available procedures for the resultants we need to calculate. Indeed we have been able to complete some investigations that were otherwise intractable.

## Appendix

$$\begin{aligned}
P_1 = & d^{11}f^{11} - 6d^{10}f^{11}m - 4d^{10}f^{10}m + 11d^{10}f^{10} + 16d^9f^{11}m^2 - 2d^9f^{11}m + \\
& 24d^9f^{10}m^2 - 54d^9f^{10}m + 4d^9f^9m^2 - 44d^9f^9m + 55d^9f^9 - 26d^8f^{11}m^3 \\
& + 12d^8f^{11}m^2 - 2d^8f^{11}m - 62d^8f^{10}m^3 + 121d^8f^{10}m^2 - 12d^8f^{10}m - \\
& 28d^8f^9m^3 + 208d^8f^9m^2 - 222d^8f^9m + 4d^8f^8m^3 + 55d^8f^8m^2 - \\
& 214d^8f^8m + 165d^8f^8 + 30d^7f^{11}m^4 - 30d^7f^{11}m^3 + 9d^7f^{11}m^2 + \\
& 90d^7f^{10}m^4 - 180d^7f^{10}m^3 + 74d^7f^{10}m^2 - 16d^7f^{10}m + 86d^7f^9m^4 - \\
& 428d^7f^9m^3 + 405d^7f^9m^2 - 24d^7f^9m - 16d^7f^8m^4 - 254d^7f^8m^3 + \\
& 804d^7f^8m^2 - 552d^7f^8m - 8d^7f^7m^4 - 4d^7f^7m^3 + 292d^7f^7m^2 - \\
& 608d^7f^7m + 330d^7f^7 - 26d^6f^{11}m^5 + 40d^6f^{11}m^4 - 16d^6f^{11}m^3 -
\end{aligned}$$

$$\begin{aligned}
& 80d^6f^{10}m^5 + 205d^6f^{10}m^4 - 176d^6f^{10}m^3 + 59d^6f^{10}m^2 - 150d^6f^9m^5 + \\
& 538d^6f^9m^4 - 542d^6f^9m^3 + 198d^6f^9m^2 - 56d^6f^9m + 16d^6f^8m^5 + \\
& 507d^6f^8m^4 - 1282d^6f^8m^3 + 767d^6f^8m^2 + 42d^6f^7m^5 + 20d^6f^7m^4 - \\
& 976d^6f^7m^3 + 1836d^6f^7m^2 - 924d^6f^7m + 2d^6f^6m^5 - 36d^6f^6m^4 - \\
& 144d^6f^6m^3 + 836d^6f^6m^2 - 1120d^6f^6m + 462d^6f^6 + 16d^5f^{11}m^6 - \\
& 30d^5f^{11}m^5 + 14d^5f^{11}m^4 + 44d^5f^{10}m^6 - 166d^5f^{10}m^5 + 204d^5f^{10}m^4 - \\
& 82d^5f^{10}m^3 + 160d^5f^9m^6 - 492d^5f^9m^5 + 633d^5f^9m^4 - 466d^5f^9m^3 + \\
& 165d^5f^9m^2 + 20d^5f^8m^6 - 40d^5f^7m^5 + 1323d^5f^7m^4 - 2136d^5f^7m^3 + \\
& 857d^5f^7m^2 + 84d^5f^7m - 14d^5f^6m^6 + 136d^5f^6m^5 + 354d^5f^6m^4 - \\
& 2140d^5f^6m^3 + 2756d^5f^6m^2 - 1092d^5f^6m + 2d^5f^5m^6 + 24d^5f^5m^5 - \\
& 8d^5f^5m^4 - 540d^5f^5m^3 + 1460d^5f^5m^2 - 1400d^5f^5m + 462d^5f^5 - \\
& 6d^4f^{11}m^7 + 12d^4f^{11}m^6 - 6d^4f^{11}m^5 - 14d^4f^{10}m^7 + 79d^4f^{10}m^6 - \\
& 116d^4f^{10}m^5 + 51d^4f^{10}m^4 - 104d^4f^9m^7 + 332d^4f^9m^6 - 522d^4f^9m^5 + \\
& 464d^4f^9m^4 - 170d^4f^9m^3 - 60d^4f^8m^7 + 517d^4f^8m^6 - 1084d^4f^8m^5 + \\
& 1112d^4f^8m^4 - 740d^4f^8m^3 + 255d^4f^8m^2 + 90d^4f^7m^7 + 52d^4f^7m^6 - \\
& 1064d^4f^7m^5 + 1642d^4f^7m^4 - 890d^4f^7m^3 + 310d^4f^7m^2 - 140d^4f^7m + \\
& 40d^4f^6m^7 - 208d^4f^6m^6 - 288d^4f^6m^5 + 1877d^4f^6m^4 - 2090d^4f^6m^3 + \\
& 501d^4f^6m^2 + 168d^4f^6m - 10d^4f^5m^7 - 78d^4f^5m^6 + 112d^4f^5m^5 + \\
& 1046d^4f^5m^4 - 2990d^4f^5m^3 + 2844d^4f^5m^2 - 924d^4f^5m - 6d^4f^4m^6 + \\
& 28d^4f^4m^5 + 178d^4f^4m^4 - 960d^4f^4m^3 + 1634d^4f^4m^2 - 1204d^4f^4m + \\
& 330d^4f^4 + d^3f^{11}m^8 - 2d^3f^{11}m^7 + d^3f^{11}m^6 + 2d^3f^{10}m^8 - \\
& 16d^3f^{10}m^7 + 26d^3f^{10}m^6 - 12d^3f^{10}m^5 + 38d^3f^9m^8 - 140d^3f^9m^7 + \\
& 235d^3f^9m^6 - 202d^3f^9m^5 + 69d^3f^9m^4 + 56d^3f^8m^8 - 318d^3f^8m^7 + \\
& 694d^3f^8m^6 - 838d^3f^8m^5 + 586d^3f^8m^4 - 180d^3f^8m^3 - 40d^3f^7m^8 - \\
& 68d^3f^7m^7 + 664d^3f^7m^6 - 1264d^3f^7m^5 + 1243d^3f^7m^4 - 770d^3f^7m^3 + \\
& 235d^3f^7m^2 - 60d^3f^6m^8 + 192d^3f^6m^7 + 120d^3f^6m^6 - 944d^3f^6m^5 + \\
& 1058d^3f^6m^4 - 476d^3f^6m^3 + 222d^3f^6m^2 - 112d^3f^6m + 20d^3f^5m^8 + \\
& 72d^3f^5m^7 - 220d^3f^5m^6 - 422d^3f^5m^5 + 1483d^3f^5m^4 - 1132d^3f^5m^3 + \\
& 31d^3f^5m^2 + 168d^3f^5m + 24d^3f^4m^7 - 114d^3f^4m^6 - 202d^3f^4m^5 + \\
& 1574d^3f^4m^4 - 2758d^3f^4m^3 + 2028d^3f^4m^2 - 552d^3f^4m - 6d^3f^3m^6 - \\
& 36d^3f^3m^5 + 357d^3f^3m^4 - 956d^3f^3m^3 + 1180d^3f^3m^2 - 704d^3f^3m + \\
& 165d^3f^3 - 6d^2f^9m^9 + 26d^2f^9m^8 - 42d^2f^9m^7 + 30d^2f^9m^6 - \\
& 8d^2f^9m^5 - 24d^2f^8m^9 + 121d^2f^8m^8 - 260d^2f^8m^7 + 294d^2f^8m^6 - \\
& 172d^2f^8m^5 + 41d^2f^8m^4 - 2d^2f^7m^9 + 76d^2f^7m^8 - 330d^2f^7m^7 + \\
& 654d^2f^7m^6 - 712d^2f^7m^5 + 414d^2f^7m^4 - 100d^2f^7m^3 + 50d^2f^6m^9 -
\end{aligned}$$

$$\begin{aligned}
& 148d^2f^6m^8 - 12d^2f^6m^7 + 557d^2f^6m^6 - 998d^2f^6m^5 + 942d^2f^6m^4 - \\
& 520d^2f^6m^3 + 129d^2f^6m^2 - 20d^2f^5m^9 + 12d^2f^5m^8 + 68d^2f^5m^7 + \\
& 8d^2f^5m^6 - 198d^2f^5m^5 + 162d^2f^5m^4 - 90d^2f^5m^3 + 114d^2f^5m^2 - \\
& 56d^2f^5m - 36d^2f^4m^8 + 174d^2f^4m^7 - 179d^2f^4m^6 - 280d^2f^4m^5 + \\
& 618d^2f^4m^4 - 238d^2f^4m^3 - 155d^2f^4m^2 + 96d^2f^4m + 18d^2f^3m^7 + \\
& 30d^2f^3m^6 - 496d^2f^3m^5 + 1350d^2f^3m^4 - 1644d^2f^3m^3 + 964d^2f^3m^2 - \\
& 222d^2f^3m + 6d^2f^2m^6 - 76d^2f^2m^5 + 295d^2f^2m^4 - 544d^2f^2m^3 + \\
& 532d^2f^2m^2 - 268d^2f^2m + 55d^2f^2 + 4df^8m^{10} - 20df^8m^9 + 40df^8m^8 - \\
& 40df^8m^7 + 20df^8m^6 - 4df^8m^5 + 8df^7m^{10} - 48df^7m^9 + 121df^7m^8 - \\
& 64df^7m^7 + 126df^7m^6 - 52df^7m^5 + 9df^7m^4 - 22df^6m^{10} + 88df^6m^9 - \\
& 90df^6m^8 - 106df^6m^7 + 334df^6m^6 - 324df^6m^5 + 146df^6m^4 - 26df^6m^3 \\
& + 10df^5m^{10} - 48df^5m^9 + 116df^5m^8 - 234df^5m^7 + 425df^5m^6 - \\
& 560df^5m^5 + 458df^5m^4 - 206df^5m^3 + 39df^5m^2 + 24df^4m^9 - 118df^4m^8 \\
& + 252df^4m^7 - 334df^4m^6 + 320df^4m^5 - 194df^4m^4 + 28df^4m^3 + 38df^4m^2 \\
& - 16df^4m - 18df^3m^8 + 48df^3m^7 + 15df^3m^6 - 150df^3m^5 + 120df^3m^4 + \\
& 48df^3m^3 - 93df^3m^2 + 30df^3m - 12df^2m^7 + 108df^2m^6 - 366df^2m^5 + \\
& 624df^2m^4 - 576df^2m^3 + 276df^2m^2 - 54df^2m + 6dfm^6 - 40dfm^5 + \\
& 111dfm^4 - 164dfm^3 + 136dfm^2 - 60dfm + 11df - 2f^7m^{11} + 12f^7m^{10} \\
& - 30f^7m^9 + 40f^7m^8 - 30f^7m^7 + 12f^7m^6 - 2f^7m^5 + 4f^6m^{11} - \\
& 24f^6m^{10} + 60f^6m^9 - 80f^6m^8 + 60f^6m^7 - 24f^6m^6 + 4f^6m^5 - 2f^5m^{11} \\
& + 18f^5m^{10} - 68f^5m^9 + 142f^5m^8 - 180f^5m^7 + 142f^5m^6 - 68f^5m^5 + \\
& 18f^5m^4 - 2f^5m^3 - 6f^4m^{10} + 30f^4m^9 - 49f^4m^8 + 105f^4m^6 - \\
& 154f^4m^5 + 105f^4m^4 - 36f^4m^3 + 5f^4m^2 + 6f^3m^9 - 42f^3m^8 + \\
& 124f^3m^7 - 198f^3m^6 + 180f^3m^5 - 86f^3m^4 + 12f^3m^3 + 6f^3m^2 - \\
& 2f^3m + 6f^2m^8 - 32f^2m^7 + 66f^2m^6 - 60f^2m^5 + 10f^2m^4 + \\
& 24f^2m^3 - 18f^2m^2 + 4f^2m - 6fm^7 + 36fm^6 - 90fm^5 + 120fm^4 - \\
& 90fm^3 + 36fm^2 - 6fm + m^6 - 6m^5 + 15m^4 - 20m^3 + 15m^2 - 6m + 1,
\end{aligned}$$

$$\begin{aligned}
P_2 = & d^9f^9 + 2d^8f^{10}m - 2d^8f^{10} - 7d^8f^9m + 4d^8f^9 - 4d^8f^8m + 7d^8f^8 - \\
& 10d^7f^{10}m^2 + 10d^7f^{10}m + 17d^7f^9m^2 + 2d^7f^9m - 16d^7f^9 + 20d^7f^8m^2 \\
& - 57d^7f^8m + 32d^7f^8 + 8d^7f^7m^2 - 27d^7f^7m + 20d^7f^7 + 20d^6f^{10}m^3 \\
& - 20d^6f^{10}m^2 - 17d^6f^9m^3 - 50d^6f^9m^2 + 66d^6f^9m - 38d^6f^8m^3 + \\
& 126d^6f^8m^2 - 30d^6f^8m - 56d^6f^8 - 36d^6f^7m^3 + 138d^6f^7m^2 - \\
& 215d^6f^7m + 112d^6f^7 - 6d^6f^6m^3 + 51d^6f^6m^2 - 73d^6f^6m + 28d^6f^6 - \\
& 20d^5f^{10}m^4 + 20d^5f^{10}m^3 + 3d^5f^9m^4 + 100d^5f^9m^3 - 103d^5f^9m^2 +
\end{aligned}$$

$$\begin{aligned}
& 32d^5f^8m^4 - 90d^5f^8m^3 - 128d^5f^8m^2 + 186d^5f^8m + 64d^5f^7m^4 - \\
& 253d^5f^7m^3 + 411d^5f^7m^2 - 110d^5f^7m - 112d^5f^7 + 26d^5f^6m^4 - \\
& 195d^5f^6m^3 + 430d^5f^6m^2 - 485d^5f^6m + 224d^5f^6 - 44d^5f^5m^3 + \\
& 125d^5f^5m^2 - 95d^5f^5m + 14d^5f^5 + 10d^4f^{10}m^5 - 10d^4f^{10}m^4 + \\
& 7d^4f^9m^5 - 80d^4f^9m^4 + 73d^4f^9m^3 - 8d^4f^8m^5 - 21d^4f^8m^4 + \\
& 244d^4f^8m^3 - 215d^4f^8m^2 - 56d^4f^7m^5 + 200d^4f^7m^4 - 194d^4f^7m^3 \\
& - 240d^4f^7m^2 + 290d^4f^7m - 44d^4f^6m^5 + 286d^4f^6m^4 - 692d^4f^6m^3 + \\
& 760d^4f^6m^2 - 170d^4f^6m - 140d^4f^6 - 2d^4f^5m^5 + 132d^4f^5m^4 - \\
& 485d^4f^5m^3 + 780d^4f^5m^2 - 705d^4f^5m + 280d^4f^5 + 2d^4f^4m^5 + \\
& 18d^4f^4m^4 - 101d^4f^4m^3 + 140d^4f^4m^2 - 45d^4f^4m - 14d^4f^4 - \\
& 2d^3f^{10}m^6 + 2d^3f^{10}m^5 - 5d^3f^9m^6 + 26d^3f^9m^5 - 21d^3f^9m^4 - \\
& 4d^3f^8m^6 + 51d^3f^8m^5 - 144d^3f^8m^4 + 97d^3f^8m^3 + 24d^3f^7m^6 - \\
& 57d^3f^7m^5 - 88d^3f^7m^4 + 351d^3f^7m^3 - 230d^3f^7m^2 + 36d^3f^6m^6 - \\
& 198d^3f^6m^5 + 388d^3f^6m^4 - 166d^3f^6m^3 - 330d^3f^6m^2 + 270d^3f^6m + \\
& 8d^3f^5m^6 - 144d^3f^5m^5 + 579d^3f^5m^4 - 1048d^3f^5m^3 + 855d^3f^5m^2 - \\
& 138d^3f^5m - 112d^3f^5 - 8d^3f^4m^6 - 40d^3f^4m^5 + 296d^3f^4m^4 - \\
& 685d^3f^4m^3 + 880d^3f^4m^2 - 667d^3f^4m + 224d^3f^4 - 4d^3f^3m^5 + \\
& 40d^3f^3m^4 - 89d^3f^3m^3 + 50d^3f^3m^2 + 31d^3f^3m - 28d^3f^3 + d^2f^9m^7 - \\
& 2d^2f^9m^6 + d^2f^9m^5 + 2d^2f^8m^7 - 16d^2f^8m^6 + 26d^2f^8m^5 - 12d^2f^8m^4 \\
& - 4d^2f^7m^7 - 2d^2f^7m^6 + 69d^2f^7m^5 - 116d^2f^7m^4 + 53d^2f^7m^3 \\
& - 14d^2f^6m^7 + 63d^2f^6m^6 - 55d^2f^6m^5 - 153d^2f^6m^4 + 289d^2f^6m^3 - \\
& 130d^2f^6m^2 - 12d^2f^5m^7 + 80d^2f^5m^6 - 233d^2f^5m^5 + 284d^2f^5m^4 + \\
& 21d^2f^5m^3 - 290d^2f^5m^2 + 150d^2f^5m + 12d^2f^4m^7 + 12d^2f^4m^6 - \\
& 246d^2f^4m^5 + 706d^2f^4m^4 - 952d^2f^4m^3 + 582d^2f^4m^2 - 58d^2f^4m - \\
& 56d^2f^4 + 12d^2f^3m^6 - 102d^2f^3m^5 + 320d^2f^3m^4 - 555d^2f^3m^3 + \\
& 610d^2f^3m^2 - 397d^2f^3m + 112d^2f^3 - 6d^2f^2m^5 + 21d^2f^2m^4 - \\
& 11d^2f^2m^3 - 37d^2f^2m^2 + 53d^2f^2m - 20d^2f^2 + df^7m^7 - 3df^7m^6 + \\
& 3df^7m^5 - df^7m^4 + 2df^6m^8 - 7df^6m^7 + 2df^6m^6 + 16df^6m^5 - 20df^6m^4 \\
& + 7df^6m^3 + 8df^5m^8 - 36df^5m^7 + 44df^5m^6 + 39df^5m^5 - 141df^5m^4 + \\
& 121df^5m^3 - 35df^5m^2 - 8df^4m^8 + 24df^4m^7 + 8df^4m^6 - 86df^4m^5 + \\
& 44df^4m^4 + 112df^4m^3 - 140df^4m^2 + 46df^4m - 12df^3m^7 + 84df^3m^6 - \\
& 275df^3m^5 + 499df^3m^4 - 491df^3m^3 + 221df^3m^2 - 10df^3m - 16df^3 + \\
& 12df^2m^6 - 61df^2m^5 + 150df^2m^4 - 236df^2m^3 + 238df^2m^2 - 135df^2m + \\
& 32df^2 - 6dfm^4 + 25dfm^3 - 39dfm^2 + 27dfm - 7df - 2f^5m^9 + 12f^5m^8 \\
& - 30f^5m^7 + 40f^5m^6 - 30f^5m^5 + 12f^5m^4 - 2f^5m^3 + 2f^4m^9 - 14f^4m^8
\end{aligned}$$



$$\begin{aligned}
& +43f^4m^7 - 75f^4m^6 + 80f^4m^5 - 52f^4m^4 + 19f^4m^3 - 3f^4m^2 + 4f^3m^8 \\
& - 22f^3m^7 + 44f^3m^6 - 30f^3m^5 - 20f^3m^4 + 46f^3m^3 - 28f^3m^2 + 6f^3m \\
& - 6f^2m^7 + 40f^2m^6 - 108f^2m^5 + 150f^2m^4 - 110f^2m^3 + 36f^2m^2 - 2f^2 \\
& - 4fm^5 + 20fm^4 - 40fm^3 + 40fm^2 - 20fm + 4f + m^5 - 5m^4 + 10m^3 \\
& - 10m^2 + 5m - 1.
\end{aligned}$$

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